



Parallel Models

Different ways to exploit parallelism

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Outline

- Shared-Variables Parallelism
 - threads
 - shared-memory architectures
- Message-Passing Parallelism
 - processes
 - distributed-memory architectures
- Practicalities
 - compilers
 - libraries
 - usage on real HPC architectures



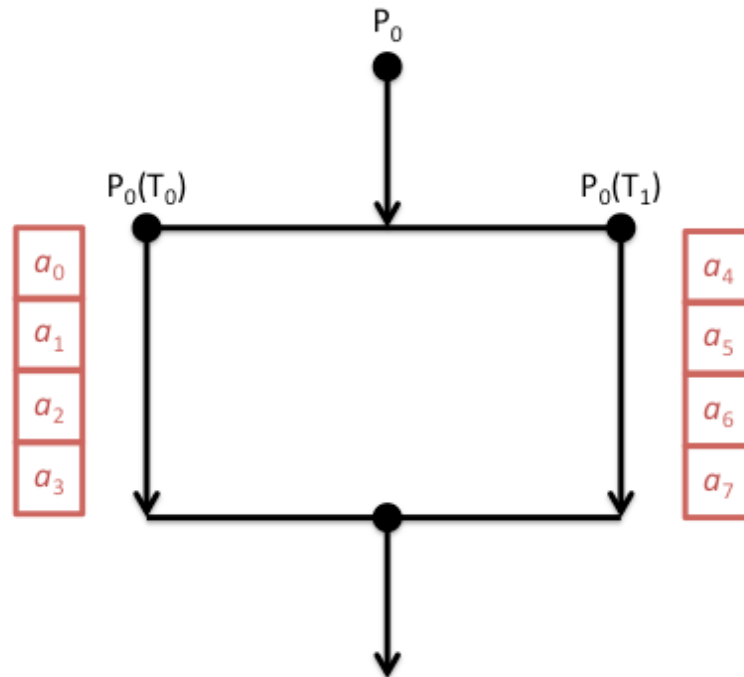
Shared Variables

Threads-based parallelism



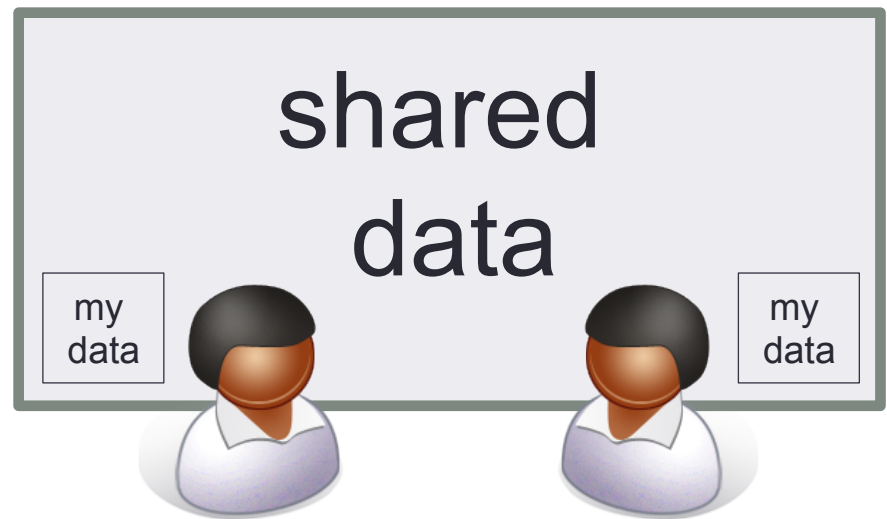
Shared-memory concepts

- Have already covered basic concepts
 - threads can all see data of parent process
 - can run on different cores
 - potential for parallel speedup



Analogy

- One very large whiteboard in a two-person office
 - the shared memory
- Two people working on the same problem
 - the threads running on different cores attached to the memory
- How do they collaborate?
 - working together
 - but not interfering
- Also need *private* data



Thread Communication

Thread 1

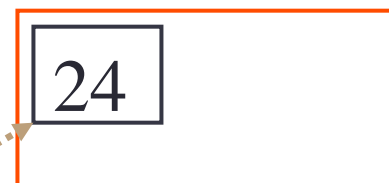
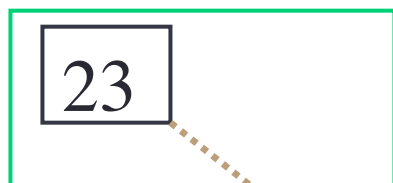
Thread 2

Program

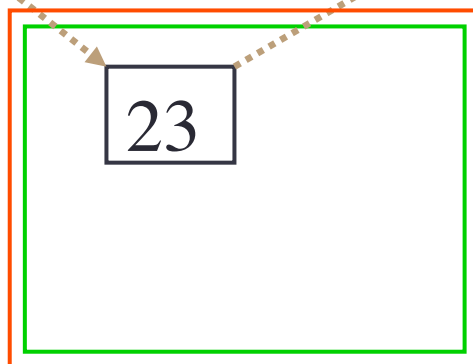
```
mya=23  
a=mya
```

```
mya=a+1
```

Private
data



Shared
data



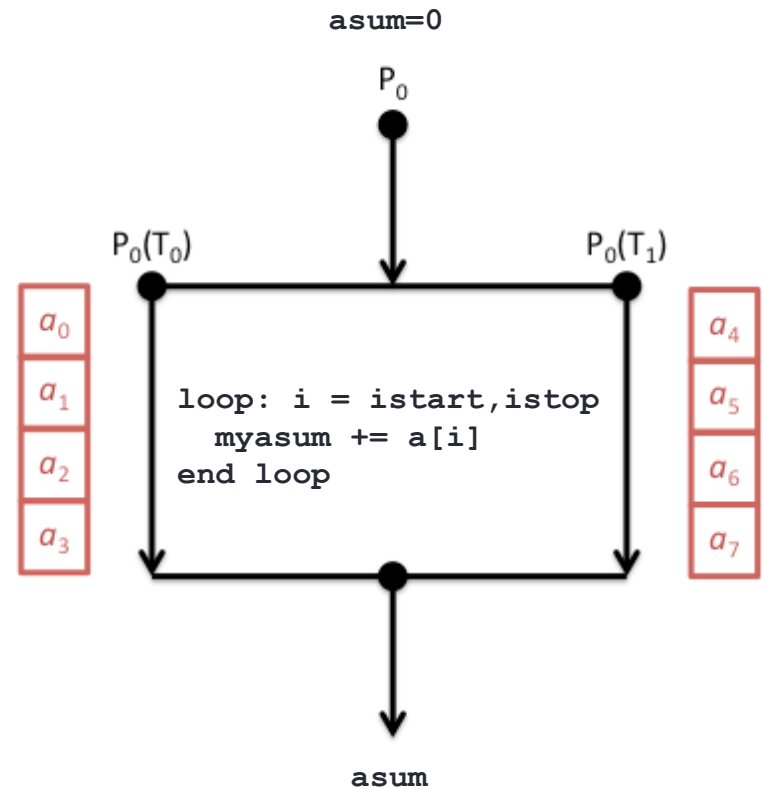
Synchronisation

- Synchronisation crucial for shared variables approach
 - thread 2's code must execute *after* thread 1
- Most commonly use global barrier synchronisation
 - other mechanisms such as locks also available
- Writing parallel codes relatively straightforward
 - access shared data as and when its needed
- Getting correct code can be difficult!



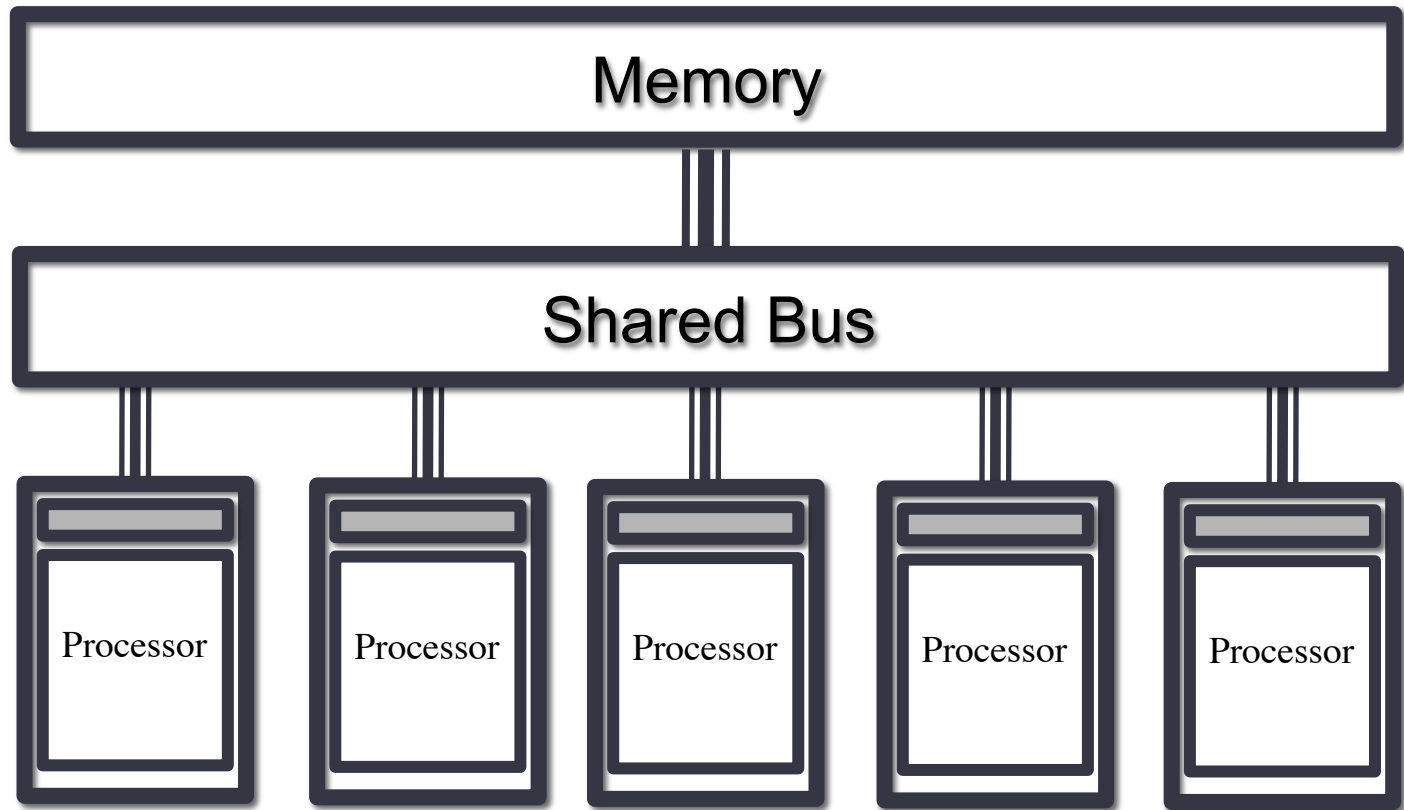
Specific example

- Computing $asum = a_0 + a_1 + \dots + a_7$
 - shared:
 - main array: `a[8]`
 - result: `asum`
 - private:
 - loop counter: `i`
 - loop limits: `istart, istop`
 - local sum: `myasum`
 - synchronisation:
 - thread0: `asum += myasum`
 - barrier
 - thread1: `asum += myasum`



Hardware

- Needs support of a shared-memory architecture



Threads: Summary

- Shared blackboard a good analogy for thread parallelism
- Requires a shared-memory architecture
 - in HPC terms, cannot scale beyond a single node
- Threads operate independently on the shared data
 - need to ensure they don't interfere; synchronisation is crucial
- Threading in HPC usually uses OpenMP directives
 - supports common parallel patterns
 - e.g. loop limits computed by the compiler
 - e.g. summing values across threads done automatically



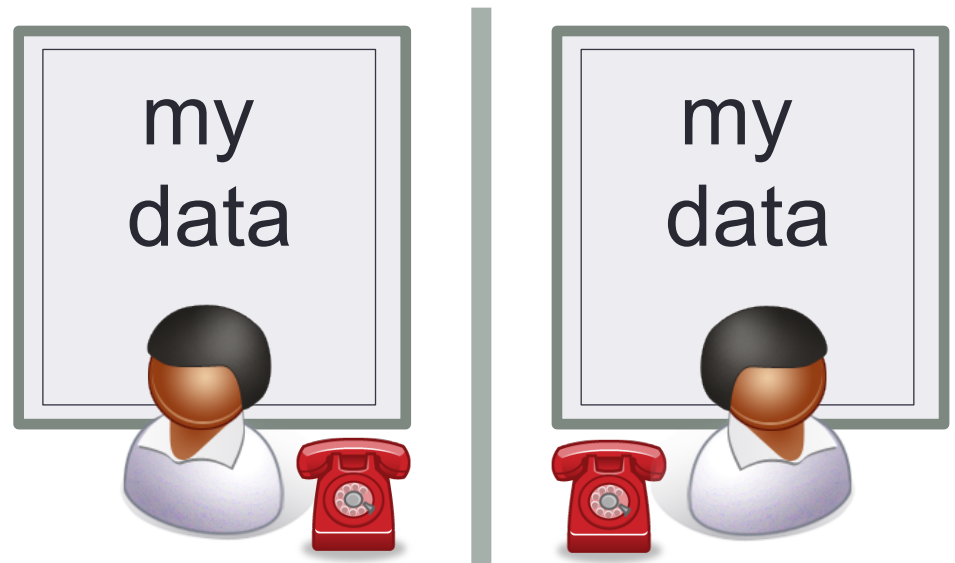
Message Passing

Process-based parallelism



Analogy

- Two whiteboards in different single-person offices
 - the distributed memory
- Two people working on the same problem
 - the processes on different nodes attached to the interconnect
- How do they collaborate?
 - to work on single problem
- Explicit communication
 - e.g. by telephone
 - no shared data



Process communication

Process 1

Program

$a=23$

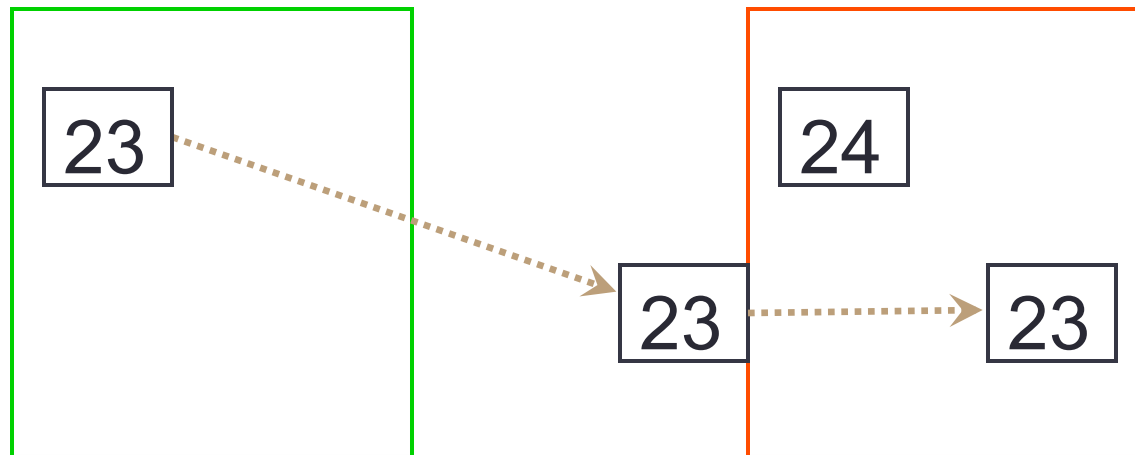
`Send(2, a)`

Process 2

`Recv(1, b)`

$a=b+1$

Data

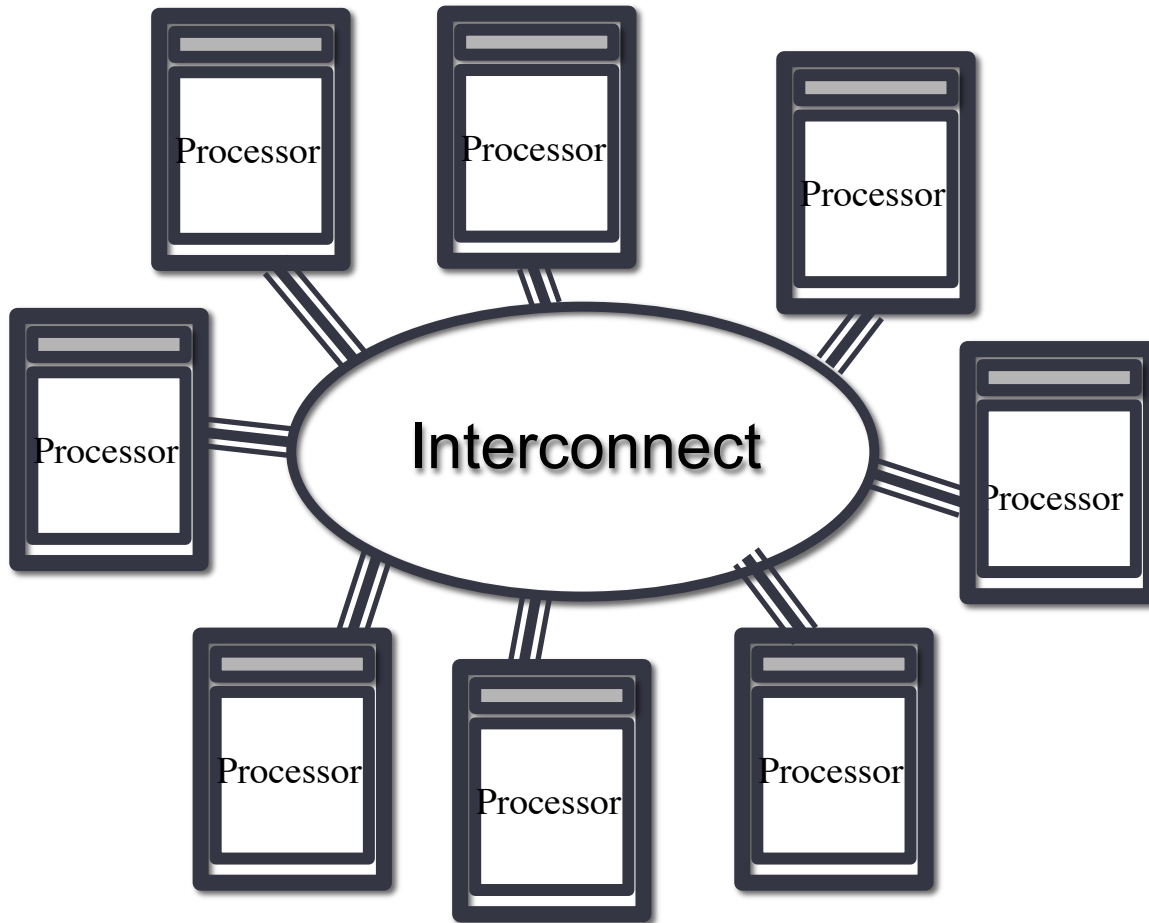


Synchronisation

- Synchronisation is automatic in message-passing
 - the messages do it for you
- Make a phone call ...
 - ... wait until the receiver picks up
- Receive a phone call
 - ... wait until the phone rings
- No danger of corrupting someone else's data
 - no shared blackboard



Hardware



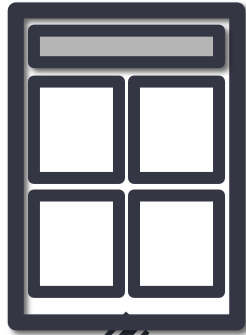
- Natural map to distributed-memory
 - one process per processor-core
 - messages go over the interconnect, between nodes/OS's

Processes: Summary

- Processes cannot share memory
 - ring-fenced from each other
 - analogous to white boards in separate offices
- Communication requires explicit *messages*
 - analogous to making a phone call, sending an email, ...
 - synchronisation is done by the messages
- Almost exclusively use Message-Passing Interface
 - MPI is a library of function calls / subroutines



Practicalities



Interconnect

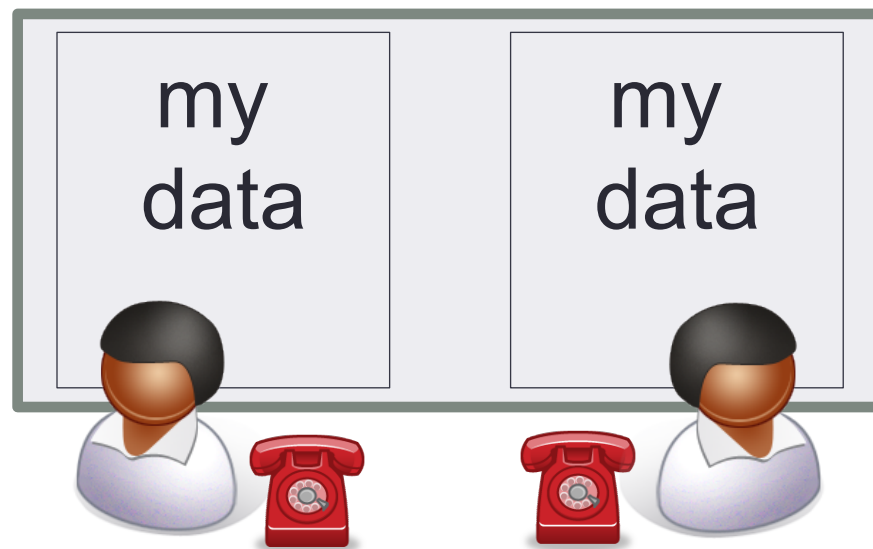


- 8-core machine might only have 2 nodes
 - how do we run MPI on a real HPC machine?
- Mostly ignore architecture
 - pretend we have single-core nodes
 - one MPI process per processor-core
 - e.g. run 8 processes on the 2 nodes
- Messages between processes on the same node are fast
 - but remember they also share access to the network

Message Passing on Shared Memory

- Run one process per core
 - don't directly exploit shared memory
 - analogy is phoning your office mate
 - actually works well in practice!

- Message-passing programs run by a special job launcher
 - user specifies #copies
 - some control over allocation to nodes



Summary

- Shared-variables parallelism
 - uses threads
 - requires shared-memory machine
 - easy to implement but limited scalability
 - in HPC, done using OpenMP compilers
- Distributed memory
 - uses processes
 - can run on any machine: messages can go over the interconnect
 - harder to implement but better scalability
 - on HPC, done using the MPI library

