



ARCHER Single Node Optimisation

Profiling

Slides contributed by Cray and EPCC





What is profiling?

- Analysing your code to find out the proportion of execution time spent in different routines.
- Essential to know this if we are going to target optimisation.
- No point optimising routines that don't significantly contribute to the overall execution time.
 - can just make your code less readable/maintainable





Code profiling

- Code profiling is the first step for anyone interested in performance optimisation
- Profiling works by instrumenting code at compile time
 - Thus it's (usually) controlled by compiler flags
 - Can reduce performance
- Standard profiles return data on:
 - Number of function calls
 - Amount of time spent in sections of code
- Also tools that will return hardware specific data
 - Cache misses, TLB misses, cache re-use, flop rate, etc...
 - Useful for in-depth performance optimisation





Sampling and tracing

- Many profilers work by sampling the program counter at regular intervals (normally 100 times per second).
 - low overhead, little effect on execution time
- Builds a statistical picture of which routines the code is spending time in.
 - if the run time is too small (< ~10 seconds) there aren't enough samples for good statistics
- Tracing can get more detailed information by recording some data (e.g. time stamp) at entry/exit to functions
 - higher overhead, more effect on runtime
 - unrestrained use can result in huge output files





Standard Unix profilers

- Standard Unix profilers are prof and gprof
- Many other profiling tools use same formats
- Usual compiler flags are -p and -pg:

```
• ftn -p mycode.F90 -o myprog for prof
```

- cc -pg mycode.c -o myprog for gprof
- When code is run it produces instrumentation log
 - mon.out for prof
 - gmon.out for gprof
- Then run prof/gprof on your executable program
 - eg. gprof myprog (not gprof gmon.out)





Standard profilers

• prof myprog reads mon.out and produces this:

%Time	Seconds	Cumsecs	#Calls	msec/call	Name					
32.4	0.71	0.71	14	50.7	relax_					
28.3	0.62	1.33	14	44.3	resid_					
11.4	0.25	1.58	3	83.	f90_close					
5.9	0.13	1.71	1629419	0.0001	_mcount					
5.0	0.11	1.82	339044	0.0003	f90_slr_i4					
5.0	0.11	1.93	167045	0.0007						
inrange_single										
2.7	0.06	1.99	507	0.12	_read					
2.7	0.06	2.05	1	60.	MAIN					





Standard profilers

- gprof myprog reads gmon.out and produces something very similar
- gprof also produces a program calltree sorted by inclusive times
- Both profilers list all routines, including obscure system ones
 - Of note: mcount(), _mcount(), moncontrol(), _moncontrol()
 monitor() and _monitor() are all overheads of the profiling
 implementation itself
 - _mcount() is called every time your code calls a function; if it's high in the profile, it can indicate high function-call overhead
 - gprof assumes calls to a routine from different parents take the same amount of time – may not be true





The Golden Rules of profiling

- Profile your code
 - The compiler/runtime will NOT do all the optimisation for you.
- Profile your code yourself
 - Don't believe what anyone tells you. They're wrong.
- Profile on the hardware you want to run on
 - Don't profile on your laptop if you plan to run on ARCHER.
- Profile your code running the full-sized problem
 - The profile will almost certainly be qualitatively different for a test case.
- Keep profiling your code as you optimise
 - · Concentrate your efforts on the thing that slows your code down.
 - This will change as you optimise.
 - So keep on profiling.





CrayPAT

 Can do both statistic sampling and function/loop level tracing.

Recommended usage:

- 1. Build and instrument code
- 2. Run code and get statistic profile
- 3. Re-instrument based on profile
- 4. Re-run code to get more detailed tracing





Example with CrayPAT (1/2)

- Load performance tools software module load perftools
- Re-build application (keep .o files)
 make clean
 make
- Instrument application for automatic profiling analysis
 - You should get an instrumented program a.out+pat pat_build -O apa a.out
- Run the instrumented application (...+pat) to get top time consuming routines
 - You should get a performance file ("<sdatafile>.xf") or multiple files in a directory <sdatadir>





Example with CrayPAT (2/2)

- Generate text report and an .apa instrumentation file pat_report [<sdatafile>.xf | <sdatadir>]
 - Inspect the .apa file and sampling report whether additional instrumentation is needed
 - See especially sites "Libraries to trace" and "HWPC group to collect"
- Instrument application for further analysis (a.out+apa)
 pat_build -0 <apafile>.apa
- Run application (...+apa)
- Generate text report and visualization file (.ap2)
 pat_report -o my_text_report.txt <data>
- View report in text and/or with Cray Apprentice² app2 <datafile>.ap2





Finding single-core hotspots

- Remember: pay attention only to user routines that consume significant portion of the total time
- View the key hardware counters, for example
 - L1 and L2 cache metrics
 - use of vector (SSE/AVX) instructions
 - Computational intensity (= ratio of floating point ops / memory accesses)

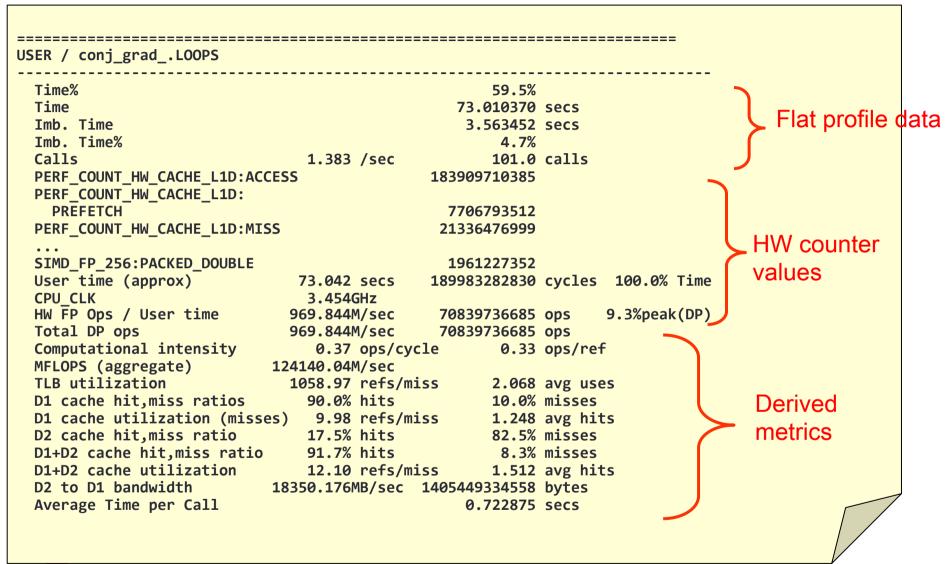




- CrayPAT has mechanisms for finding "the" hotspot in a routine (e.g. in case the routine contains several and/or long loops)
 - CrayPAT API
 - Possibility to give labels to "PAT regions"
 - Loop statistics (works only with Cray compiler)
 - Compile & link with CCE using -h profile_generate
 - pat_report will generate loop statistics if the flag is enabled











Hardware performance counters

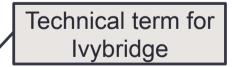
- CrayPAT can interface with Cray XC30's HWPCs
 - Gives extra information on how hardware is behaving
 - Very useful for understanding (& optimising) application performance
- Provides information on
 - hardware features, e.g. caches, vectorisation and memory bandwidth
- Available on per-program and per-function basis
 - Per-function information only available through tracing
- Number of simultaneous counters limited by hardware
 - 4 counters available with Intel Ivybridge processors
 - If you need more, you'll need multiple runs
- Most counters accessed through the PAPI interface
 - Either native counters or derived metrics constructed from these





Hardware counters selection

- HWPCs collected using CrayPAT
 - Compile and instrument code for profiling as before
- Set PAT RT HWPC environment variable at runtime
 - e.g. in the job script
 - Hardware counter events are not collected by default (except with APA)
- export PAT_RT_HWPC=...
 - either a list of named PAPI counters
 - or <set number> = a pre-defined (and useful) set of counters
 - recommended way to use HWPCs
 - there are around 20 groups to choose from
 - To see them:
 - pat_help -> counters -> ivybridge -> groups
 - man hwpc
 - more \${CRAYPAT_ROOT}/share/CounterGroups.intel_fam6mod62







Predefined Ivybridge HW Counter Groups

Default is number 1 with CrayPAT APA procedure

0: D1 with instruction counts

1: Summary -- FP and cache metrics

2: D1, D2, L3 Metrics

6: Micro-op queue stalls

7: Back end stalls

8: Instructions and branches

9: Instruction cache

10: Cache Hierarchy

11: Floating point operations dispatched

12: AVX floating point operations

13: SSE and AVX floating point

operations SP

14: SSE and AVX floating point

operations DP

19: Prefetchs

23: FP and cache metrics (same as 1)





Time%		1.6%		
Time		0.909054	secs	
Imb. Time		0.057555	secs	
Imb. Time%		6.4%		
Calls	0.116M/sec	187500.0	calls	
PAPI_L1_DCM	18.108M/sec	29376518	misses	
PAPI_TLB_DM	0.007M/sec	11643	misses	
PAPI_L1_DCA	170.243M/sec	276182686	refs	
PAPI_FP_OPS		0	ops	
DATA_CACHE_REFILLS_FROM_L2_OR_NORTHE	BRIDGE:			
ALL	18.711M/sec	30354680	fills	
DATA_CACHE_REFILLS_FROM_NORTHBRIDGE	0.003M/sec	5084	fills	
User time (approx)	1.622 secs	3731260602	cycles	100.0% Time
HW FP Ops / User time		0	ops	0.0%peak(DP)
HW FP Ops / WCT				
Computational intensity	0.00 ops/c	ycle 0.00	ops/ref	:
MFLOPS (aggregate)	0.00M/sec			
TLB utilization	23720.03 refs/	miss 46.328	avg use	2S
D1 cache hit,miss ratios	89.4% hits	10.6%	misses	
D1 cache hit,refill ratio	89.0% hits	11.0%	refills	5
D1 cache utilization (misses)	9.40 refs/		avg hit	
	9.10 refs/		avg use	25
D2 cache hit, miss ratio	100.0% hits			
D1+D2 cache hit, miss ratio	100.0% hits	0.0%	misses	

Some hints on interpreting the data

TLB utilization

- Memory loaded in pages: 4kB standard (could use larger hugepages)
- e.g. 512 x 8-byte double precision floats
- So if every double was used once, expect 512 refs/miss
 - Less than 512 shows poor use; more than 512 is good (5420.38 excellent)
 - N.B. <avg uses> = <refs/miss> / 512





D1 cache utilization

- Level 1 cache line is 64 contiguous bytes, e.g. 8 x 8-byte doubles
- So if every double was used once, expect 8 refs/miss
 - Corresponds to hit ratio of 87.5% [i.e. 100*(1 1/<refs/miss>)]
 - N.B. <avg uses> = <refs/miss> / 8
 - Less than 8 (or 87.5%) shows poor use
 - Rule of thumb: want this to be 20 (or 95%) or more

D1+D2 cache hit ratio

Should be high (rule of thumb is more than 97%);





CrayPAT observations and suggestions

D1 + D2 cache utilization: 39.8% of total execution time was spent in 4 functions with combined D1 and D2 cache hit ratios below the desirable minimum of 97.0%. Cache utilization might be improved by modifying the alignment or stride of references to data arrays in these functions.

```
D1_D2_cache_hit_ratio Time% Function

56.8% 12.0% calc3_

77.9% 6.4% calc2_

95.7% 1.4% calc1_

96.3% 20.0% calc3_.LOOP@li.80
```

TLB utilization: 19.6% of total execution time was spent in 3 functions with fewer than the desirable minimum of 512 data references per TLB miss. TLB utilization might be improved by modifying the alignment or stride of references to data arrays in these functions.



Interpreting the performance numbers

- Performance numbers are an average over all ranks
 - explains non-integer values
- This does not always make sense
 - e.g. if ranks are not all doing the same thing:
 - Master-slave schemes
 - MPMD apruns combining multiple, different programs
- Want them to only process data for certain ranks
 - pat_report -sfilter_input='condition' ...
 - condition should be an expression involving pe, e.g.
 - pe<1024 for the first 1024 ranks only
 - pe%2==0 for every second rank





OpenMP data collection and reporting

- Give finer-grained profiling of threaded routines
 - Measure overhead incurred entering and leaving
 - Parallel regions
 - #pragma omp parallel
 - Work-sharing constructs within parallel regions
 - #pragma omp for
- Timings and other data now shown per-thread
 - rather than per-rank
- OpenMP tracing enabled with pat_build -gomp ...
 - CCE: insert tracing points around parallel regions automatically
 - Intel, Gnu: need to use CrayPAT API manually





OpenMP data collection and reporting

- Load imbalance for hybrid MPI/OpenMP programs
 - now calculated across all threads in all ranks
 - imbalances for MPI and OpenMP combined
 - Can choose to see imbalance in each programming model separately
 - See next slide for details
- Data displayed by default in pat_report
 - no additional options needed
 - Report focuses on where program is spending its time
 - Assumes all requested resources should be used
 - you may have reasons not to want to do this, of course





Imbalance options for data display (pat_report -0 ...)

- These options control how load balance is displayed:
- profile_pe_th (default view)
 - Imbalance based on the set of all threads in the program
 - · i.e. imbalance from OpenMP and MPI combined
 - this is best measure to understand code performance
- profile pe.th
 - Highlights imbalance across MPI ranks
 - Thread data for each rank is aggregated
 - · max used rather than mean, to highlight under-performers
 - Aggregated thread data merged into MPI rank data
- profile_th_pe
 - For each thread, show imbalance over MPI ranks
 - Example: Load imbalance shows where thread 4 in each MPI rank didn't get much work





Memory usage

- Knowing how much memory each rank uses is important:
 - What is the minimum number of cores I can run this problem on?
 - given there is 32GB (~30GB usable) of memory per node (32 cores)
 - Does memory usage scale well in the application?
 - Is memory usage balanced across the ranks in the application?
 - Is my application spending too much time allocating and freeing?
- Profile heap usage using CrayPAT
 - pat_build -gheap ...





Heap statistics,

Memory per rank

~30GB usable memory per node

Notes for table 5:

Too many allocs/frees? Would show up as ETC time in CrayPAT report

Table option:

-O heap_hiwater

Options implied by table option:
-d am@,ub,ta,ua,tf,nf,ac,ab -b pe=[mmm]

This table shows only lines with Tracked Heap HiWater MBytes > 0.

Heap Stats during Main Program Table 5:

	Tracked Heap HiWater MBytes	/otal Allocs	Total Frees	Tracked Objects Not Freed	Tracked MBytes Not Freed	PE[mmm]
1	9.794	915	910	4	1.011	Total
	9.943 9.909 9.446	1170 715 1278	1103 712 1275	68 3 3	1.046 1.010 1.010	pe.0 pe.22 pe.43

Memory leaks

Not usually a problem in HPC

Summary

- Profiling is essential to identify performance bottlenecks
 - even at single core level
- CrayPAT has some very useful extra features
 - can pinpoint and characterise the hotspot loops (not just routines)
 - hardware performance counters give extra insight into performance
 - well-integrated view of hybrid programming models
 - most commonly MPI/OpenMP
 - also CAF, UPC, SHMEM, pthreads, OpenACC, CUDA
 - information on memory usage
- And remember the Golden Rules
 - including the one about not believing what anyone tells you



