# Performance of Parallel IO on Lustre and GPFS

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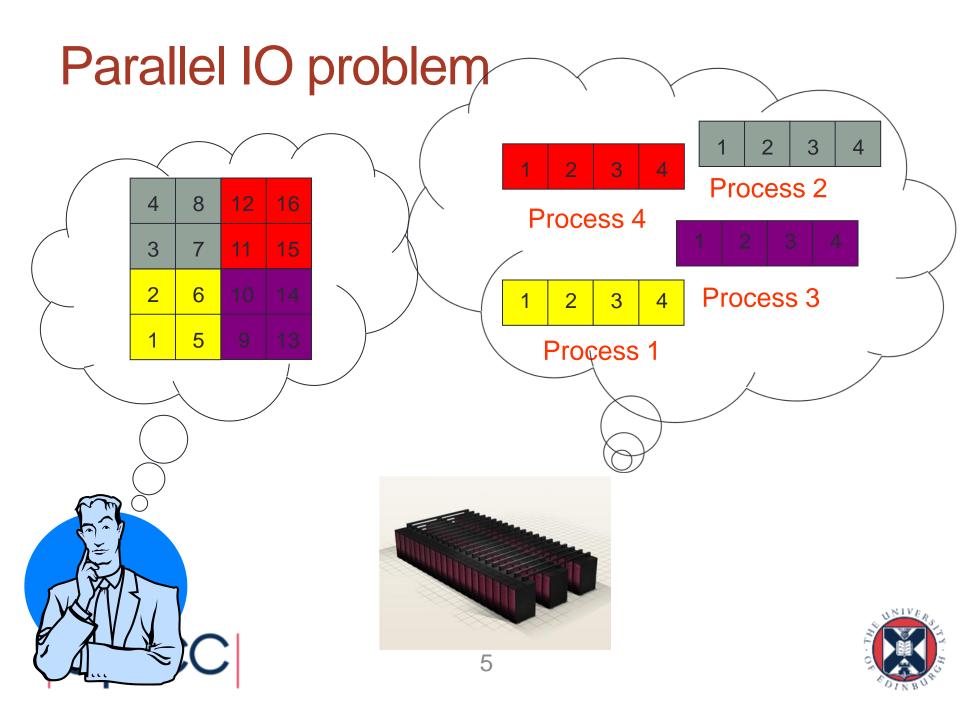




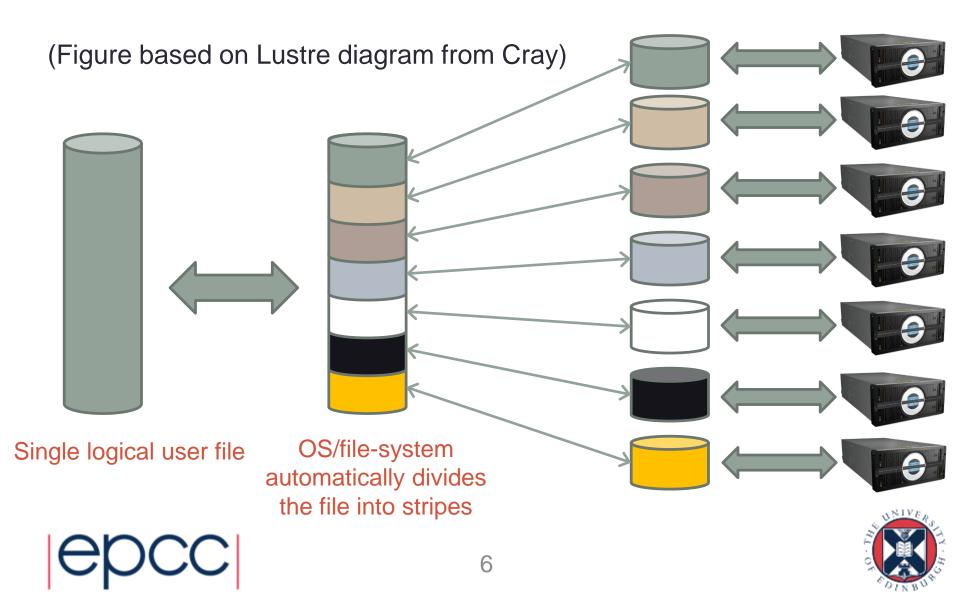
# Outline

- Parallel IO problem
- Common IO patterns
- Parallel filesystems
- MPI-IO Benchmark results
- Filesystem tuning
- MPI-IO Application results
- HDF5 and NetCDF
- Conclusions





## **Parallel Filesystems**



# Common IO patterns

- Multiple files, multiple writers
  - each process writes its own file
  - numerous usability and performance issues
- Single file, single writer (master IO)
  - high usability but poor performance
- Single file, multiple writers
  - all processes write to a single file; poor performance
- Single file, collective writers
  - aggregate data onto a subset of IO processes
  - hard to program and may require tuning
  - potential for scalable IO performance





# **Quantifying Performance**

- What is good performance on ARCHER?
  - Generally see ~500MB/s per OST
  - This is the serial limit. If getting that, not achieving parallel I/O
- Always benchmark and quantify bandwidth
  - Use the Cray performance tools
- Contention is an issue can see huge variance in results
  - Do multiple runs at different times of day
  - Look at best and worst case
- Beware of caching effects on performance



#### Performance – Large Number of Files

*"setting striping to 1 has reduced total read time for his 36000 small files from 2 hours to 6 minutes"* 

- comment on resolution of an ARCHER helpdesk query.
- User was performing I/O on 36000 separate files of ~300KB with 10000 processes
- Had set parallel striping to maximum possible (48 OSTs / -1) assuming this would give best performance
- Overhead of querying every OST for every file dominated the access time
- Moral: more stripes does not mean better performance





### Performance – Large Number of Files 2

- 15GB consisting of 5500 1.5-4MB files
  - Effect of striping on serial "tar" operation:

\$> time tar -cf stripe48.tar stripe48

real 31m19.438s

\$> time tar -cf stripe4.tar stripe4
real 24m50.604s

\$> time tar -cf stripe1.tar stripe1
real 18m34.475s

~40% reduction in operation time between 48 and 1 stripe

- Still bottlenecks at MDS. This access pattern is not recommended, but it is common.

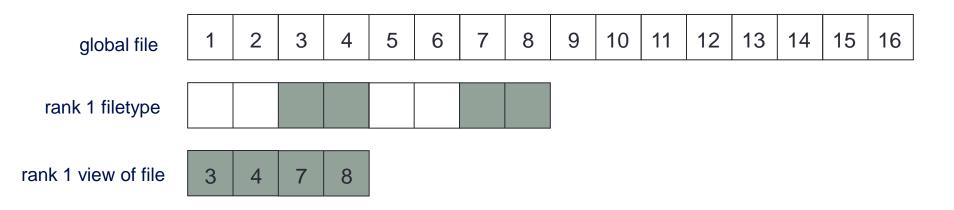




## **Global description: MPI-IO**



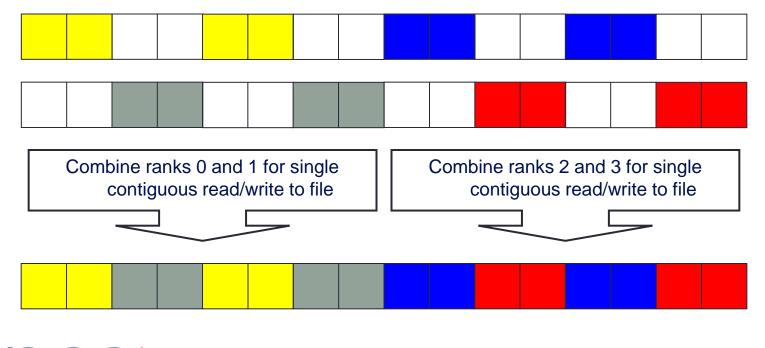
rank 1	rank 3
(0,1)	(1,1)
rank 0	rank 2
(0,0)	(1,0)





# Collective IO

- Enables numerous optimisations in principle
  - requires global description and participation of all processes
  - does this help in practice?





## Cellular Automaton Model

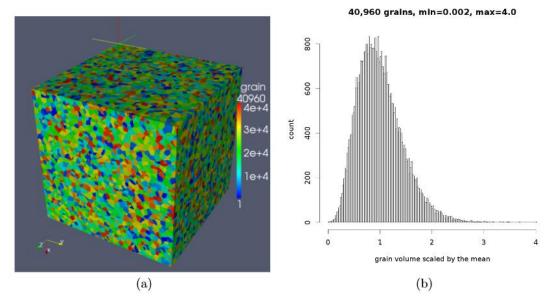


Figure 1: A  $4.1 \times 10^9$  cell, 40,960 grain equiaxed microstructure model, showing (a) grain arrangement with colour denoting orientation; (b) grain size size (volume) histogram.

 Fortran coarray library for 3D cellular automata microstructure simulation, Anton Shterenlikht, proceedings of 7<sup>th</sup> International Conference on PGAS Programming Models, 3-4 October 2013, Edinburgh, UK. 13



## Benchmark

Distributed regular 3D dataset across 3D process grid

- local data has halos of depth 1; set up for weak scaling
- implemented in Fortran and MPI-IO

! After opening file **fh**, define what portions of file this process owns call **MPI\_File\_set\_view(fh**, disp, MPI\_DOUBLE\_PRECISION, **filetype**,

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'native', MPI\_INFO\_NULL, ierr)

! Write data collectively

call MPI\_File\_write\_all(fh, iodata, 1, mpi\_subarray, status, ierr)



#### **ARCHER XC30**







# EPSRC

Engineering and Physical Sciences Research Council







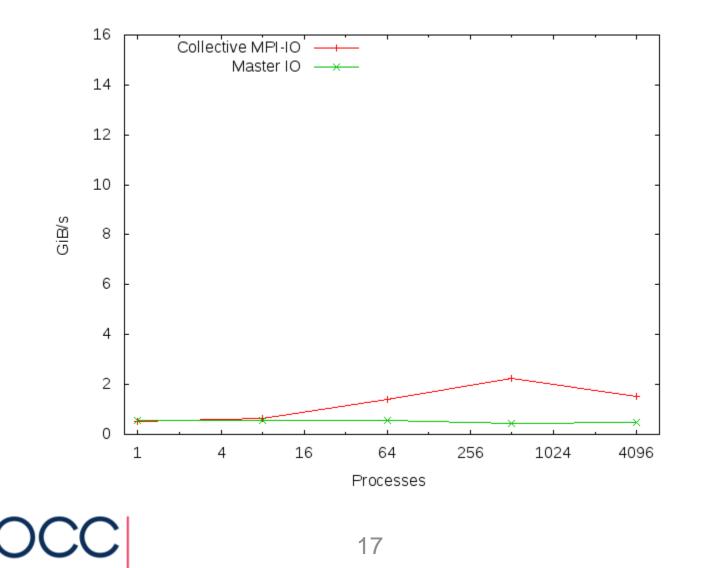
# Single file, multiple writers

- Serial bandwidth on ARCHER around 400 to 500 MiB/s
- Use MPI\_File\_write not MPI\_File\_write\_all
  - identical functionality
  - different performance

Processes	Bandwidth
1	49.5 MiB/s
8	5.9 MiB/s
64	2.4 MiB/s



## Single file, collective writers



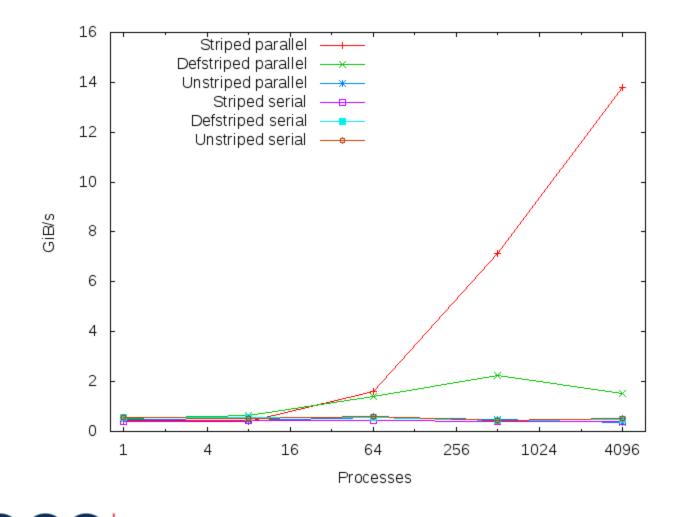


# Lustre striping

- We've done a lot of work to enable (many) collective writers
  - learned MPI-IO and described data layout to MPI
  - enabled collective IO
  - MPI dynamically decided on number of writers
  - collected data and aggregates before writing
- ... for almost no benefit!
- Need many physical disks as well as many IO streams
  - in Lustre, controlled by the number of stripes
  - default number of stripes is 4; ARCHER has around 50 IO servers
- User needs to set striping count on a per-file/directory basis
  - lfs setstripe -c -1 <directory> # use maximal striping

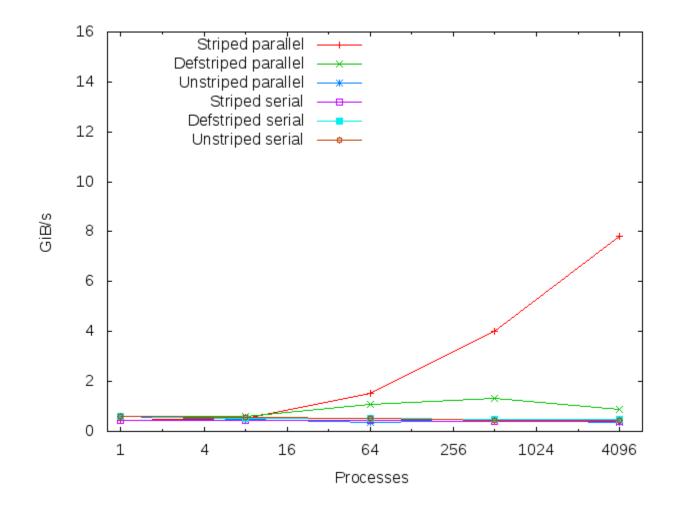


## Cray XC30 with Lustre: 128<sup>3</sup> per proc





## Cray XC30 with Lustre: 256<sup>3</sup> per proc



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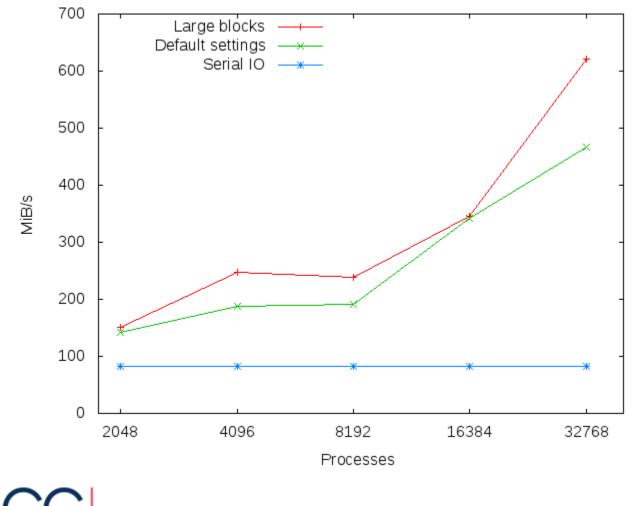


# Performance Summary

- Serial IO never gets more than about 500 MiB/s
  - peak for a single OST
- With default striping, never exceed 2 GiB/s
  - 4 stripes = 4 OSTs = 4 x 500 MiB/s
- With full striping, IO bandwith increases with process count
  - can achieve in excess of 10 GiB/s
- Collective IO is essential
  - replacing MPI\_File\_Write\_all() by MPI\_File\_write()
    disastrous!
  - identical functionality but each IO request now processed separately with file locking



### BG/Q: #IO servers scales with CPUs





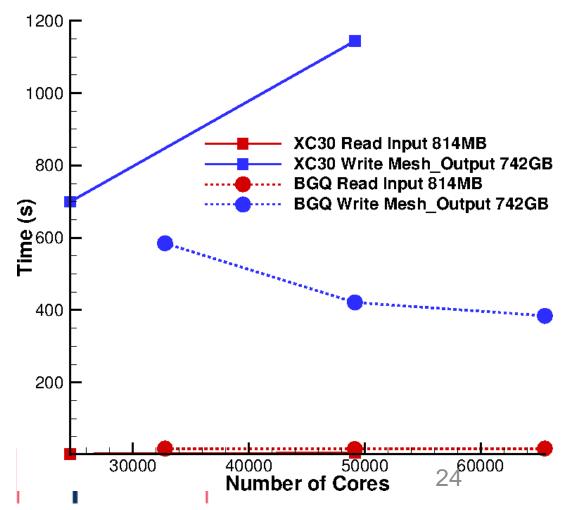
## Code\_Saturne http://code-saturne.org

- CFD code developed by EDF (France)
- Co-located finite volume, arbitrary unstructured meshes, predictor-corrector
- 350 000 lines of code
  - 50% C
  - 37% Fortran
  - 13% Python
- MPI for distributed-memory (some OpenMP for sharedmemory) including MPI-IO
- Laminar and turbulent flows: k-eps, k-omega, SST, v2f, RSM, LES models, ... EPCC 23



## Code\_SATURNE: default settings

MPI-IO - 7.2 B Tetra Mesh

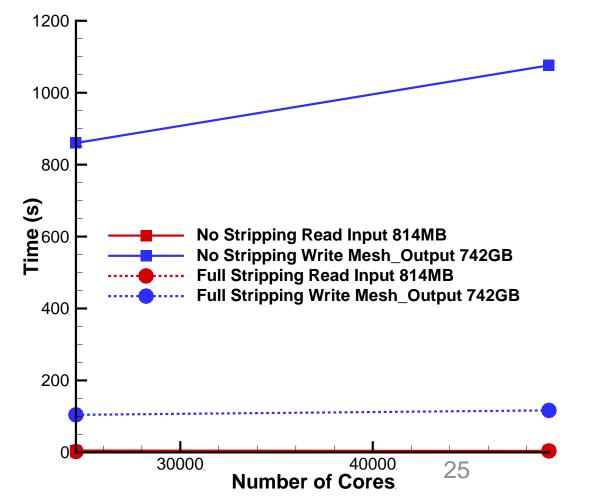


- Consistent with benchmark results
  - default striping Lustre similar to GPFS



## Code\_Saturne: Lustre striping

MPI-IO - 7.2 B Tetra Mesh

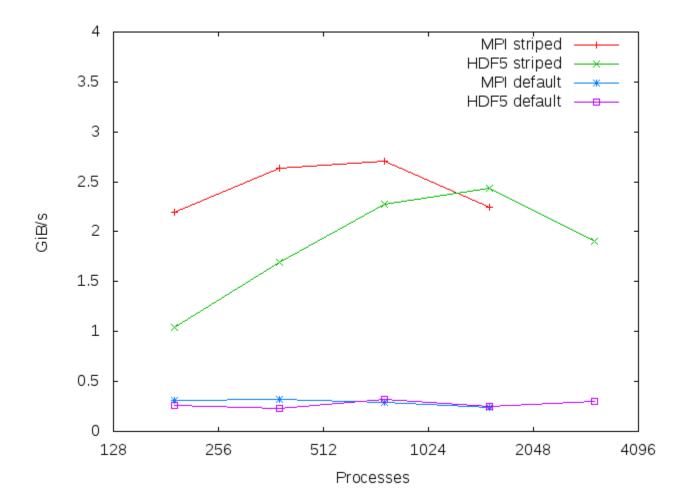


Consistent with benchmark results

 order of magnitude improvement from striping



## Simple HDF5 benchmark: Lustre



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# Further Work

- Non-blocking parallel IO could hide much of writing time
  - or use more restricted split-collective functions
  - extend benchmark to overlap comms with calculation
- I don't believe it is implemented in current MPI-IO libraries
   blocking MPI collectives are used internally
- A subset of user MPI processes will be used by MPI-IO
  - would be nice to exclude them from calculation
  - extend MPI\_Comm\_split\_type() to include something like MPI\_COMM\_TYPE\_IONODE as well as MPI\_COMM\_TYPE\_SHARED ?



# Conclusions

- Efficient parallel IO requires all of the following
  - a global approach
  - coordination of multiple IO streams to the same file
  - collective writers
  - filesystem tuning
- MPI-IO Benchmark useful to inform real applications
  - NetCDF and HDF5 layered on top of MPI-IO
  - although real application IO behaviour is complicated
- Try a library before implementing bespoke solutions!
  - higher level view pays dividends

